Corrections (1997)

p. 19 (Fig. 1.2.5) Change A_3 to A^3

p. 22 (Fig. 1.2.8) Change top $4, 1$ to $4, 0$

p. 42 (Fig. 1.3.4) Change top $\sum_{i=1}^{5} a_i$ to $\sum_{i=1}^{9} a_i$

p. 65 (-9) Change "any single" to "an optimal single"

p. 67 (Table 3.3) The upper bound $p - 1$ for the total exchange has been reduced to $p/2$; see the paper "Optimal Communication Algorithms for Hypercubes," J. of Parallel and Distributed Computation, Vol. 11, 1991, pp. 263-275, by D. P. Bertsekas, C. Ozveren, G. Stamoulis, P. Tseng, and J. N. Tsitsiklis.

p. 80 (+15) Change "at every time ..." to "at the first time unit, each node sends its packet to all its neighbors. Afterwards, at every time ..."

p. 86 (+7) Change "bo" to "be"

p. 86 (+9) Change $2 \log n$ to $4 \log n$

p. 86 (-1) Change "2 time units using the links illustrated" to "4 time units using the transfers illustrated"

p. 88 (+5) Change "processors." to "processors arranged in a square array."

p. 103 (-9) Change "that" to "than"

p. 106 (+13) Change ρ_S^{t-D} to ρ_S^t

p. 130 (+17) Change $L'DL$ to LDL'

p. 130 (+19) Delete "square"

p. 179 (+5) Change "[Ash]" to "[Ash70]"

p. 240 (Proof of part (c)) This proof works for the case where the penalty parameter sequence $c(t)$ is bounded above. If $c(t) \to \infty$ a slightly different argument is needed.

p. 264 (+16) Change λ_j to λ_i

p. 271 (+5) Change 3.13 to 3.1.3

p. 307 (-2) Change $\min\{x_{i(k+1)}^k, x_{(k+1)j}^k\}$ to $\max\{x_{i(k+1)}^k, x_{(k+1)j}^k\}$

p. 327 (+13) Change $[x \mid x \geq 0, x_1 = 0]$ to $\{x \mid x \geq 0, x_1 = 0\}$

p. 335 (+17) Change "capacity" to "demand"

p. 339 (-3) Change p_i to p_j

p. 340 (+10) Change ":=" to "="; change p_i to p_j

p. 363 (+2) Change b_{ij} to b_{ji}

p. 368 (-3) Change “assigment” to “assignment”

p. 375 (+6) Change “$\epsilon < 1/n$” to “$\epsilon < 1/m$, where m is the number of similarity classes”

p. 375 (-9) Change “Exercise 3.7” to “Exercise 3.6”

p. 386 (+16) Change $O(|N|^3+$ to $O(M|N|^3+$

p. 390 (Exercise 4.5) Change all occurances of “8” to “6”

p. 392 (Figure 5.5.1) Change “Slope $= b_{ij}$” to “Slope $= -b_{ij}$”; change “Slope $= c_{ij}$” to “Slope $= -c_{ij}$”

p. 396 (-3) In the middle figure change $q(t_{ij})$ to $q_{ij}(t_{ij})$

p. 386 (+16) Change $O(|N|^3+$ to $O(M|N|^3+$

p. 408 (-9) Change $f_{ij} \leq 0$ to $f_{ij} < 0$

p. 423 (-7) Change “result” to “results”

p. 429 (-7,-8,-12,-13) Change all X to x

p. 430 (-10) Change $\tau_4^2(9)$ to $\tau_3^2(9)$

p. 459 (+3) Change $R_i(p)$ to R_i

p. 548 (+12) Change “rerouting” to “routing”

p. 569 (+6) Change “of a large” to “for a large”

p. 571 For a recent treatment of the problem of termination detection, including additional methods and references, see the paper “Finite Termination of Asynchronous Iterative Algorithms,” Parallel Computing, Vol. 22, 1996, pp. 39-56, by S. A. Savari and D. P. Bertsekas.

p. 618 (+3) Change [GaB87] to [GaB86]

p. 620 (+8) Change $\ln n$ to $\ln x$

p. 620 (+17) Delete “$< x, y >$”

p. 649 (-4) Include as an additional assumption that Z is closed

p. 659 (-8) Change “A set C” to “A nonempty set C”

p. 659 (-1) A proof of the closure of the cone C appears in many sources, including “Nonlinear Programming,” by D. P. Bertsekas, Athena Scientific, 1995, p. 580.